# MACM 401/MATH 701/MATH 801 Assignment 4, Spring 2017.

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This assignment is to be handed in by Monday March 13th by 4pm. For problems involving Maple calculations and Maple programming, you should submit a printout of a Maple worksheet of your Maple session. Late Penalty: -20% for up to 48 hours late. Zero after that.

## Question 1: *P*-adic Lifting (20 marks)

Reference: Section 6.2 and 6.3.

- (a) By hand, determine the *p*-adic representation of the integer u = 116 for p = 5, first using the positive representation, then using the symmetric representation for  $\mathbb{Z}_5$ .
- (b) Theorem: Let  $u, p \in \mathbb{Z}$  with p > 2. If  $-\frac{p^n}{2} \le u < \frac{p^n}{2}$  there exist unique integers  $u_0, u_1, \ldots, u_{n-1}$  such that  $u = u_0 + u_1 p + \cdots + u_{n-1} p^{n-1}$  and  $-\frac{p}{2} \le u_i < \frac{p}{2}$ .

Prove uniqueness.

(c) Determine the cube-root, *if it exists*, of the following polynomials

$$a(x) = x^{6} - 531x^{5} + 94137x^{4} - 5598333x^{3} + 4706850x^{2} - 1327500x + 125000,$$
  

$$b(x) = x^{6} - 406x^{5} + 94262x^{4} - 5598208x^{3} + 4706975x^{2} - 1327375x + 125125$$

using reduction mod 5 and linear *p*-adic lifting. You will need to derivive the update formula by modifying the update formula for computing the  $\sqrt{a(x)}$ .

Factor the polynomials so you know what the answers are. Express your the answer in the p-adic representation. To calculate the initial solution  $u_0 = \sqrt[3]{a} \mod 5$  use any method. Use Maple to do all the calculations.

#### Question 2: Hensel lifting (15 marks)

Reference: Section 6.4 and 6.5.

(a) Given

$$a(x) = x^4 - 2x^3 - 233x^2 - 214x + 85$$

and image polynomials

$$u_0(x) = x^2 - 3x - 2$$
 and  $w_0(x) = x^2 + x + 3$ ,

satisfying  $a \equiv u_0 w_0 \pmod{7}$ , lift the image polynomials using Hensel lifting to find (if there exist) u and w in  $\mathbb{Z}[x]$  such that a = uw.

(b) Given

$$b(x) = 48 x^4 - 22 x^3 + 47 x^2 + 144$$

and an image polynomials

$$u_0(x) = x^2 + 4x + 2$$
 and  $w_0 = x^2 + 4x + 5$ 

satisfying  $b \equiv 6 u_0 w_0 \pmod{7}$ , lift the image polynomials using Hensel lifting to find (if there exist) u and w in  $\mathbb{Z}[x]$  such that b = uw.

### Question 3: Determinants (25 marks)

Consider the following 3 by 3 matrix A of polynomials in  $\mathbb{Z}[x]$  and its determinant d.

> P := () -> randpoly(x,degree=2,dense): > A := Matrix(3,3,P);

$$A := \begin{bmatrix} -55 - 7x^2 + 22x & -56 - 94x^2 + 87x & 97 - 62x \\ -83 - 73x^2 - 4x & -82 - 10x^2 + 62x & 71 + 80x^2 - 44x \\ -10 - 17x^2 - 75x & 42 - 7x^2 - 40x & 75 - 50x^2 + 23x \end{bmatrix}$$

> d := LinearAlgebra[Determinant](A);

 $d := -224262 - 455486 x^2 + 55203 x - 539985 x^4 + 937816 x^3 + 463520 x^6 - 75964 x^5$ 

(a) (15 marks) Let A by an n by n matrix of polynomials in  $\mathbb{Z}[x]$  and let  $d = \det(A)$ . Develop a modular algorithm for computing  $d = \det(A) \in \mathbb{Z}[x]$ . Your algorithm will compute determinants of A modulo a sequence of primes and apply the CRT. For each prime p it will compute the determinant in  $\mathbb{Z}_p[x]$  by evaluation and interpolation. In this way we reduce computation of a determinant of a matrix over  $\mathbb{Z}[x]$  to many computations of determinants of matrices over  $\mathbb{Z}_p$ , a field, for which ordinary Gaussian elimination, which does  $O(n^3)$  arithmetic operations in  $\mathbb{Z}_p$ , may be used.

You will need bounds for deg d and  $||d||_{\infty}$ . Use primes p = [101, 103, 107, ...] and use Maple to do Chinese remaindering. Use x = 1, 2, 3, ... for the evaluation points and use Maple for interpolation. Implement your algorithm in Maple and test it on the above example.

To reduce the coefficients of the polynomials in A modulo p = 7 in Maple use

 $> B := A \mod p;$ 

To evaluate the polynomials in B at  $x = \alpha$  modulo p in Maple use

> C := eval(B,x=alpha) mod p;

To compute the determinant of a matrix C over  $\mathbb{Z}_p$  in Maple use

> Det(C) mod p;

(b) (10 marks) Suppose A is an n by n matrix over  $\mathbb{Z}[x]$  and  $A_{i,j} = \sum_{k=0}^{d} a_{i,j,k} x^k$  and  $|a_{i,j,k}| < B^m$ . That is A is an n by n matrix of polynomials of degree at most d with coefficients at most m base B digits long. Assume the primes satisfy  $B and that arithmetic in <math>\mathbb{Z}_p$  costs O(1). Estimate the time complexity of your algorithm in big O notation as a function of n, m and d. Make reasonable simplifying assumptions such as n < B and d < B as necessary. State your assumptions. Also helpful is

$$\ln n! < n \ln n \quad \text{for} \quad n > 1.$$

#### Question 4: A linear *x*-adic Newton iteration (15 marks).

Let p be an odd prime and let  $a(x) = a_0 + a_1 x + ... + a_n x^n \in \mathbb{Z}_p[x]$  with  $a_0 \neq 0$  and  $a_n \neq 0$ . Suppose  $\sqrt{a_0} = \pm u_0 \mod p$ . The goal of this question is to design an x-adic Newton iteration algorithm that given  $u_0$ , determines if  $u = \sqrt{a(x)} \in \mathbb{Z}_p[x]$  and if so computes u.

(a) Let

$$u = u_0 + u_1 x + \dots + u_{k-1} x^{k-1} + u_k x^k + \dots$$

Derive the Newton update formula for  $u_k$ . Show your working.

(b) Now test your update formula on the two polynomials  $a_1(x)$  and  $a_2(x)$  below using p = 101 and  $u_0 = +5$ . Please print out the sequence of values of  $u_0, u_1, u_2, ...$  as you compute them. Note, one of the polynomials has a  $\sqrt{in \mathbb{Z}_p[x]}$ , the other does not. So you will need to work out when the algorithm should stop lifting. Do all calculations in Maple.

 $a_1 = 81 x^6 + 16 x^5 + 24 x^4 + 89 x^3 + 72 x^2 + 41 x + 25$  $a_2 = 81 x^6 + 46 x^5 + 34 x^4 + 19 x^3 + 72 x^2 + 41 x + 25$ 

(c) The update formula requires  $u_0 \neq 0$ . Explain briefly what you should you do if  $a_0 = 0$  and you want to compute  $\sqrt{a(x)} \in \mathbb{Z}_p[x]$  if it exists.