

Computing in  $\mathbb{Q}(\alpha_1, \dots, \alpha_n)$ .

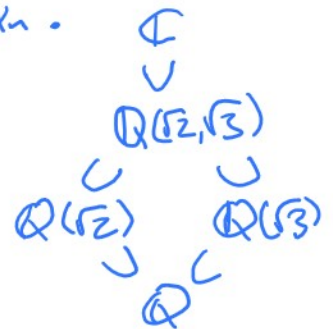
Def. Let  $\alpha_1, \dots, \alpha_n$  be algebraic numbers.

$\mathbb{Q}(\alpha_1, \dots, \alpha_n)$  is the smallest field containing  $\mathbb{Q}, \alpha_1, \dots, \alpha_n$ .

So  $\mathbb{Q}(\alpha_1, \dots, \alpha_n) \subset \mathbb{C}$ .

E.g.  $\mathbb{Q}(\sqrt{2}, \sqrt{3})$ .

Notice  $\mathbb{Q}(\sqrt{2}), \mathbb{Q}(\sqrt{3})$  are subfields of  $\mathbb{Q}(\sqrt{2}, \sqrt{3})$ .



How can we compute in  $\mathbb{Q}(\alpha_1, \dots, \alpha_n)$ ?

Use  $\mathbb{Q}(\alpha_1, \dots, \alpha_n) \cong \mathbb{Q}(\alpha_1)(\alpha_2)(\alpha_3) \dots (\alpha_n)$ .

$$L_0 := \mathbb{Q}$$

$$L_1 := \mathbb{Q}(\alpha_1) \cong \mathbb{Q}[z_1]/m_1(z_1) \text{ where } m_1 \text{ is the min. poly for } \alpha_1 \text{ over } \mathbb{Q}.$$

$$L_2 := L_1[z_2]/m_2(z_2) \text{ where } m_2 \text{ " " " " " } \alpha_2 \text{ over } L_1.$$

⋮

$$L_n := L_{n-1}[z_n]/m_n(z_n) \text{ where } m_n \text{ " " " " " } \alpha_n \text{ over } L_{n-1}.$$

Let  $d_k = \deg(m_k, z_k) \geq 1$ .  $L_k$  is a quotient ring.

$\Rightarrow L_k \cong L_{k-1}^{\overset{d_k \leftarrow \dim(L_k)}{\text{as a vector space with basis}}}$

$$B_k = \{1, z_k, z_k^2, \dots, z_k^{d_k-1}\}.$$

Since  $m_k$  is irreducible over  $L_{k-1}$ ,  $L_k$  is a field and

$$L_k \cong \mathbb{Q}(\alpha_1, \dots, \alpha_k) \text{ for } 1 \leq k \leq n.$$

E.g.  $\alpha_1 = \sqrt{2}, \alpha_2 = \sqrt{3}$ .

$$z_1 = \sqrt{2} \Rightarrow z_1^2 = 2 \Rightarrow \boxed{z_1^2 - 2} = 0 \Rightarrow m_1 = z_1^2 - 2.$$

$$\Rightarrow L_1 = \mathbb{Q}[z_1]/(z_1^2 - 2), \quad B = \{1, z_1\}, \quad d_1 = 2.$$

Notice  $\sqrt{3} \notin \mathbb{Q}(\sqrt{2})$ . so  $\deg(m_2, z_2) > 1$ .

Try:  $m_2(z_2) = 1 \cdot z_2^2 + (a_1 z_1 + a_0) z_2 + (b_1 z_1 + b_0) \in L_1[z_2]$ .

$0 = m_2(\sqrt{3}) = 3 + (a_1 z_1 + a_0) z_2 + (b_1 z_1 + b_0)$

$\Rightarrow 0 = a_1 z_1 z_2 + a_0 z_2 + b_1 z_1 + (3 + b_0)$

$\Rightarrow a_1 = 0, a_0 = 0, b_1 = 0, b_0 = -3$

$\Rightarrow m_2(z_2) = z_2^2 - 3. \quad B_2 = \{1, z_2\}$

$L_2 = \{ [a_1 + a_2 z_2] : a_1, a_2 \in L_1 \}$

We could represent  $[(3 + 2z_1) \cdot 1 + (7 + 5z_1) \cdot z_2] \in L_2$   
as  $[[3, 2], [7, 5]]$ . I do this in recden.

E.g.  $\alpha_1 = \sqrt{2}, \alpha_2 = 1 + \sqrt{2} + \sqrt{3}$ .

$m_1(z_1) = z_1^2 - 2$

$m_2(z_2) = z_2^2 - (2z_1 + 2)z_2 + z_1 z_2 \in L_1[z_2]$

We also have  $L_n \cong \mathbb{Q}[z_1, \dots, z_n] / \langle m_1, m_2, \dots, m_n \rangle = R$ .

In  $\succ$ lex with  $z_1 < z_2 < \dots < z_n$   $LM(m_i(z_i)) = z_i^{d_i}$

Since  $\gcd(LM(m_i), LM(m_j)) = 1 \Rightarrow G = \langle m_1, \dots, m_n \rangle$  is a GB for  $I$ .

$R \cong \mathbb{Q}^d$  where  $d = \prod_{k=1}^n d_k$ .  $d$  is called the degree of

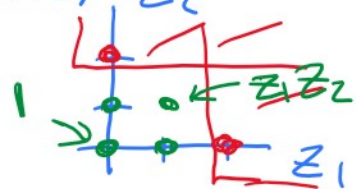
$L_n$  over  $\mathbb{Q}$  or  $\mathbb{Q}(\alpha_1, \dots, \alpha_n)$  over  $\mathbb{Q}$ . So  $R \cong \mathbb{Q}^d$ .

A basis for  $R$  is  $\overline{\langle LT(I) \rangle}$ .

E.g.  $m_1 = z_1^2 - 2, m_2 = z_2^2 - 3, I = \langle m_1, m_2 \rangle, z_2$

$\langle LT(I) \rangle = \langle z_1^2, z_2^2 \rangle$

$B = \{ 1, z_1, z_2, z_1 z_2 \}$



We could represent elements of  $R$  as multivariate polynomials in  $\mathbb{Q}[z_1, \dots, z_n]$ . Then

$[a] \cdot [b] = [a b \text{ mod } G]$ . This is quite slow.