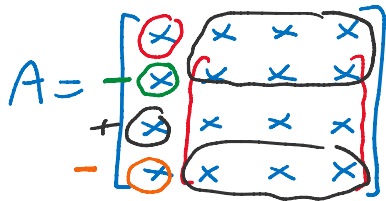


Computational Exact Linear Algebra.

Let A be an $n \times n$ matrix over a comm. ring R .
How fast can we compute $\det(A)$?

Def (cofactor expansion)



Let $B_{ij} = A$ with row i and col j removed.

$$\det(A) = \sum_{i=1}^n (-1)^{i+1} A_{i1} \cdot \det(B_{i1})$$

$$\det([a]) = a.$$

Observe $\det(A) \in R$.

Let $M(n) = \# \text{mults in } R \text{ that this method does.}$

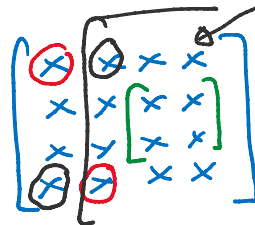
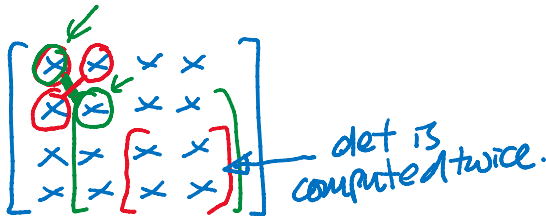
$$M(1) = 0.$$

$$M(n) = n \cdot M(n-1) + n$$

$$M(2) = 2 \cdot 0 + 2.$$

$$M(n) = 1.72 \cdot n! \text{ mults in } R.$$

Gentleman & Johnson (1976)



$\binom{4}{2}$ 2×2 submatrices from the last two columns.

① Compute all $\binom{n}{2}$ 2×2 determinants from columns $n, n-1$, using the n 1×1 determinants from column n .

② Compute all $\binom{n}{3}$ 3×3 determinants from the last 3 columns. (picking 3 rows from n) using the dets from step 1.

Repeat this until we compute all $\binom{n}{n}=1$ dets of all A .

Let $G(n) = \# \text{mults in } R \text{ that the G&J method does.}$

$$G(n) = \sum_{i=2}^n \binom{n}{i} \cdot i = \frac{1}{2} n \cdot 2^n - n \in O(n \cdot 2^n) \text{ which is exponential in } n.$$

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↑ ↑ i multiplies each
$i \times i$ minors from last i cols

$$\frac{M(10)}{G(10)} = \frac{6,255,300}{5,110} = 1220.$$

Can we compute $\det(A)$ in $O(n^3)$ ring ops, $+$, $-$, \times ?

Open.

Berkowitz (1984) $O(n^4)$

Kaltofen (1992) $O(n^{3.5})$.

What if R is an integral domain? E.g. \mathbb{Z} , $\mathbb{Z}[x_1, \dots, x_n]$.

We could run G.E. in the fraction field. $\mathbb{Q} \downarrow \mathbb{Z}(x_1, \dots, x_n)$.

$$A = \begin{bmatrix} 2 & 1 & 3 \\ 3 & -1 & 1 \\ 5 & 3 & 1 \end{bmatrix} \xrightarrow{\substack{R_2 \leftarrow R_2 - \frac{3}{2}R_1 \\ R_3 \leftarrow R_3 - \frac{5}{2}R_1}} \begin{bmatrix} 2 & 1 & 3 \\ 0 & -\frac{5}{2} & -\frac{7}{2} \\ 0 & \frac{1}{2} & -\frac{13}{2} \end{bmatrix} \xrightarrow{\substack{R_3 \leftarrow R_3 - \frac{1}{2}R_2 \\ \frac{1}{2}R_2 = +\frac{1}{5}R_2}} \begin{bmatrix} 2 & 1 & 3 \\ 0 & -\frac{5}{2} & -\frac{7}{2} \\ 0 & 0 & -\frac{36}{5} \end{bmatrix}$$

$$\text{So } \det(A) = \cancel{2} \cdot \cancel{\frac{7}{2}} \cdot \frac{-36}{5} = +36.$$

Can we avoid fractions?

$$A = \begin{bmatrix} 2 & 1 & 3 \\ 3 & -1 & 1 \\ 5 & 3 & 1 \end{bmatrix} \xrightarrow{\substack{R_2 \leftarrow 2R_2 \\ R_3 \leftarrow 2R_3}} \begin{bmatrix} 2 & 1 & 3 \\ 6 & -2 & 2 \\ 10 & 6 & 2 \end{bmatrix} \xrightarrow{\substack{R_2 \leftarrow R_2 - 3R_1 \\ R_3 \leftarrow R_3 - 5R_1}} \begin{bmatrix} 2 & 1 & 3 \\ 0 & -5 & -7 \\ 0 & 1 & -13 \end{bmatrix}$$

$$\xrightarrow{R_3 \leftarrow -5R_3} \begin{bmatrix} 2 & 1 & 3 \\ 0 & -5 & -7 \\ 0 & -5 & 65 \end{bmatrix} \xrightarrow{R_3 \leftarrow R_3 - 1 \cdot R_2} \begin{bmatrix} 2 & 1 & 3 \\ 0 & -5 & -7 \\ 0 & 0 & 72 \end{bmatrix}$$

$$\det(E) = \det(I) = -5 \det(C) = -5 \det(B) = -5 \cdot 4 \cdot \det(A).$$

$$\det(E) = 2 \cdot (-5) \cdot 72.$$

$$\det(A) = \det(E) / (-20) = \frac{2(-5) \cdot 72}{-5 \cdot 4} = \frac{72}{2} = 36.$$

The only divisions are exact because $\det(A) \in \mathbb{Z}$.

-5.4

The only divisions are exact because $\det(A) \in \mathbb{Z}$.
 But the size of the integers grows exponentially in n .
 How big is $\det(A)$? where $A_{ij} \in \mathbb{Z}$.

Hadamard's bound: $|\det(A)| \leq \prod_{i=1}^n \sqrt{\sum_{j=1}^n A_{ij}^2}$

If $|A_{ij}| < B^m$ then $|\det(A)| < \prod_{i=1}^n \sqrt{n \cdot B^{2m}} = n^{n/2} \cdot B^{nm}$

The size of $\det(A)$ $\leq \lceil \log_B B^{nm} n^{n/2} \rceil =$
 $= nm + \lceil \frac{n}{2} \log_B n \rceil \ll (m+1)n$ digits base B .
 $B=2^{64}$