

Rational Reconstruction (Wang 1981)

Suppose $\frac{n}{d} \in \mathbb{Q}$, $\gcd(n,d)=1$, $d > 0 \Rightarrow$ uniqueness.

Suppose we have computed $u = \frac{n}{d} \pmod m$, where $0 \leq u < m$,
and $\gcd(d,m)=1$. Context: $m = p_1 p_2 p_3 \dots$ or $m = p^k$.

How can we recover n/d from u, m ?

E.g. $m = 5 \cdot 7 = 35$ $\frac{n}{d} = -\frac{2}{3}$ $u = -2 \cdot 12 = -24 = +11 \pmod{35}$

How big does m need to be?

Can we recover $\frac{114}{109}$ from $\frac{114}{109} \pmod{35} = 11$?

$\Rightarrow m > 2 \cdot 114 \cdot 109$.

Run EEA with input $M, u \geq 0$.

$r_0, r_1 \leftarrow m, u$ $s_0, s_1 \leftarrow 1, 0$ $t_0, t_1 \leftarrow 0, 1$

$i \leftarrow 1$.

While $r_i \neq 0$ do

$q_{i+1} \leftarrow \lfloor \frac{r_{i-1}}{r_i} \rfloor$

$r_{i+1} \leftarrow r_{i-1} - q_{i+1} r_i$

$s_{i+1} \leftarrow s_{i-1} - q_{i+1} s_i$

$t_{i+1} \leftarrow t_{i-1} - q_{i+1} t_i$

end while

$N \leftarrow i-1$.

// $r_N = \gcd(r_0, r_1)$. $r_{N+1} = 0$.

The integers r_i, s_i and t_i satisfy

$$s_i m + t_i u = r_i \quad \text{for } 0 \leq i \leq N+1.$$

(mod m)

$$t_i \cdot u \equiv r_i \pmod{m}$$

$$\gcd(m, t_i) = 1 \Rightarrow u \equiv \underline{r_i / t_i} \pmod{m}$$

$i=0$
 $t_0=0$

$i=N+1$
 $t_{N+1}=m$
 $s_{N+1}=u$

i.e. the EEA gives us a sequence of rationals $r_i / t_i \equiv u \pmod{m}$.

Is $r_i / t_i = \frac{n}{d}$ for some $0 < i < N+1$?

Yes provided $m > |2nd|$ and $\gcd(d,m)=1$.

Which index i ?

Thm. (C. Davignon, Wang), 1982.

Which means ...

Theorem (Euy, Davenport, Wang). 1982.

Let $n, d \in \mathbb{Z}$, $d > 0$, $\gcd(n, d) = 1$.

Let $m \in \mathbb{Z}$, $m > 0$, $\gcd(m, d) = 1$ and $u = \frac{n}{d} \bmod m$ with $0 \leq u < m$.

Let $N \geq |n|$ and $D \geq d$. Then

(i) if $m > 2ND$ then $u \pmod m$ is unique in \mathbb{Z}_m . $2^{-1} = 7$.

$m=13$
 $N=3$
 $D=2$
 $2 \cdot N \cdot D = 12$

$\frac{n}{d}$	$\frac{0}{1}$	$\frac{1}{1}$	$\frac{-1}{1}$	$\frac{2}{1}$	$\frac{-2}{1}$	$\frac{3}{1}$	$\frac{-3}{1}$	$\frac{4}{2}$	$\frac{-1}{2}$	$\frac{3}{2}$	$\frac{-3}{2}$
u	0	1	12	2	10	3	9	7	6	8	5

$\frac{2}{3} \bmod 13 = 5$

(ii) if $m > 2ND$ then on input d, m, u , there exists a unique index i in EEA s.t. $r_i/t_i = \frac{n}{d}$. Moreover i is the first index s.t. $r_i \leq N$.

If we have good bounds $N \geq |n|$ and $D \geq d$ e.g. $N = 10n$ and $D = 10d$. then compute $m = p^k$ until $m > 2ND$. and apply (ii).

If we don't have good bounds? \downarrow
 E.g. solve $Ax = b$ where $\begin{bmatrix} \square & \square \\ \square & \square \end{bmatrix} \begin{bmatrix} i/2 \\ i/3 \end{bmatrix} = \begin{bmatrix} \square \\ \square \end{bmatrix}$

Wang: let $N = D = \lfloor \sqrt{m/2} \rfloor$.
 Try (ii) to get y
 Check if $Ay = b$.

EEA Cost.
 $O((\log m)^2)$.

Maximal Quotient Rational Reconstruction. Monagan 2004.

Algorithm. Output $\frac{r_i}{t_i}$ with q_{i+1} maximal.

Lemma 1.

$$\frac{m}{3} < q_{i+1} |t_i| r_i < \underline{m} \quad \forall 1 \leq i \leq N.$$

If $m \gg 2|n|d$ then q_{i+1} must be large.

If $m \gg 2|n|d$ then q_i must be large.

Maple: $\text{iratrecon}(u, m) \rightarrow \text{FAIL}$ or n/d .

$\text{iratrecon}\left(\begin{bmatrix} x \\ x \\ x \\ x \end{bmatrix}, m\right)$ $N=D=\sqrt{\frac{m}{2}}$
default

$\rightarrow \bullet x^3 + \bullet x^2 + \bullet xy + \bullet$
 $\uparrow \quad \uparrow \quad \uparrow \quad \uparrow$

$\text{iratrecon}(u, m, \text{maxquo} = \underline{1000000})$.